| Q1 | Q2 | Q3 | Q4 | Tot |
| :--- | :--- | :--- | :--- | :--- |
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## CEng 334-Introduction to Operating Systems <br> Spring 2018-2019, Midterm II, May 6, 2019

( 5 pages, 4 questions, 100 points, 100 minutes)

## METU Honor Code and Pledge

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I have read and understood the implications of the METU Honor Code. To be precise, I understand that this is a closed notes exam, and I am forbidden to access any other source of information other than provided within the exam. I will TURN OFF all my electronic equipment (phones, smart watches, etc.) and put it off the table along with other notes and materials that I may have with me. I understand that leaving electronic devices on during the exam is strictly forbidden. I understand and accept to obey all the rules announced by the course staff, and that failure to obey these will result in disciplinary action.

## Name:

$\qquad$ No: $\qquad$ Signature: $\qquad$
QUESTION 1.(25 points)
a) The kernel is not involved in any thread activity in User Level threads. These type of threads cannot utilize more than one processors
b) In a typical priority inversion scenario, there are processes $A, B$, and $C$ in increasing order of priorities (C has the highest priority). A owns a lock and $C$ waits for that lock. As a result $B$ will have a practicaly higher priority than C .
c) Which scheduling policy is non-deterministic, may create different schedules for the same input? Lottery scheduling
d) Bankers algorithm has an advantage over Resource Allocation Graph since:

> Multiple instances of resources can be handled in bankers algorithm
e) In Copy on Write, a page fault is raised due to a protection error. Page fault handler classifies the error as CoW since:

The page is marked as writable in VM area/segment but page table entry is marked as readonly
f) In case of page eviction, what happens in the following type of pages.

| page type | operation before eviction |
| :--- | :--- |
| Readonly mapped binary executable | nothing |
| Heap area modified | write to swap space |
| An memory mapped writable file, modified | write to file |
| Uninitialized area, only readonly access since allocation | nothing |

g) What type of advantage/economy does the items below provide in virtual memory? Give short answers such as 'shares memory among threads', 'keeps only one copy on disk'.

| Demand paging | only used areas loaded in memory |
| :--- | :--- |
| Copy on Write | parent and child shares pages |
| Memory mapped programs | only one copy is kept in memory |
| Translation Lookaside Buffer | fast address translation |
|  |  |

h) Most filesystems relies on the caching of metadata in the main memory, so they are accessed fast. What is the disadvantage of this?

A system crash will have metadata not stored on disk causing integrity problems during the next boot
i) Compare the efficiency of the following operations on a FAT vs indexed structure filesystem. Mark the better one with a tick $\checkmark$. Assume that the metadata is in memory. Consider only memory based operations with no disk operation.

| Operation | Indexed | FAT |
| :--- | :---: | :---: |
| Accessing a random block in a file | $\checkmark$ |  |
| Sequentially accessing all blocks of a file |  | $\checkmark$ |
| Seeking to end of file | $\checkmark$ |  |
| Inserting a new block to the beginning of a file, shifting file one block <br> right |  | $\checkmark$ |
| Deleting the file as a whole, marking all blocks as free |  | $\checkmark$ |

$\qquad$

QUESTION 2.(25 points)
Assume that you are given a 30 bits virtual memory architecture with 1024 bytes pages. PTE is 4 bytes having only 20 bits used for frame addressing. Use power of 2 notation whenever convenient.
a) Largest physical memory that can be addressed is $\qquad$ bytes
b) If one level page tables are used, A page table requires $2^{22}$ bytes.
c) If two level page tables are used with second level page table fits and uses a full page, first level can have multiple pages or shorter than a page:
Second level address is 8 bits.
First level address is 12 bits.
First level page table is $2^{14}$ bytes, fits in 16 page/s.
d) In a demand paging scenario with two level page table configuration above, draw page table after the the following full memory accesses:
$0 x 50228 a \rightarrow 0 x 068 a$,
$0 \times 504200 \rightarrow 0 \times 0 a 00$,
$0 \times 000240 \rightarrow 0 \times 1240$.
Assuming values at right handside is the pyhsical memory bytes that the given VM bytes map.
Show only valid entries in page tables as:

| index1 | frame no |
| :--- | :--- |
| index2 | frame no |

In first level pages, only draw connectors, no frame number value required.

Level One: Level Two:

| index | frameno | $\Gamma$ | index | frameno |
| :---: | :---: | :---: | :---: | :---: |
| 0x00 |  |  | 0x00 | 0x4 |
| 0x14 | $\bullet$ |  |  |  |
|  |  |  | index | frameno |
|  |  |  | 0x08 | 0x01 |
|  |  |  | 0x10 | 0x02 |

e) Assuming all PTE are valid in page tables, for one level and two levels. Which one uses more memory, and what is the difference:

Two levels uses more memory, it is 16 pages/ 16 Kilobytes more than one level
$\qquad$
$\qquad$

## QUESTION 3.(25 points)

Consider three processes which would run on the CPU as follows if they were run as the only process:

| Process $\backslash \mathrm{t}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | $\cdot$ | . | R | S | R | S | R | S | R | S | R | S | R | S | R | S | R |  |


| Process $\backslash \mathrm{t}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| B | R | R | R | S | S | S | S | S | R | R | R | S | S | S | S | S | R |  |


| Process $\backslash \mathrm{t}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| C | . | R | R | R | R | R | S | S | R | R | R | R | S | S | R | R | R |  |

where., $\mathbf{R}$ and $\mathbf{S}$ indicate that the process has not arrived yet, is $\mathbf{R u n n i n g}$ and is Sleeping respectively. Note that Process A, B and C arrives at $t=2, t=0$ and $t=1$ respectively.

Fill in the following tables, if all three processes are scheduled using the different scheduling policies. If a process is in the ready list, but is not running fill its status as $\mathbf{W}$.

- Round Robin with a CPU quantum of 3 time units.

- Shortest Job First.

- Shortest Remaining Job First.

| Process $\backslash \mathrm{t}$ | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| :--- | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| A | $\cdot$ | $\bullet$ | $\omega$ | $R$ | 4 | $R$ | $S$ | $R$ | $S$ | $R$ | $S$ | $W$ | $R$ | $S$ | $R$ | $S$ | $W$ | $R$ |
| $B$ | $R$ | $R$ | $R$ | $S$ | $S$ | $S$ | $S$ | $S$ | $R$ | $W$ | $R$ | $R$ | $S$ | $S$ | $S$ | $S$ | $S$ | $W$ |
| C | $\cdot$ | $W$ | $W$ | $W$ | $R$ | $W$ | $R$ | $W$ | $W$ | $W$ | $W$ | $W$ | $W$ | $R$ | $W$ | $R$ | $R$ | $S$ |

In both RR and SRTF, there are equally correct alternate sequences.
$\qquad$
$\qquad$

## QUESTION 4.(25 points)

Assume you have a disk with 1024 cluster/sector of $1 \mathrm{Kbytes}, 1 \mathrm{MB}$ in total and is formatted as a FAT filesystem. Sector 0 is used for boot. The consecutive sectors contain the File Allocation Table of the filesystem, followed by the root directory, which fits into one sector.
a) How many entries are there in the FAT?
b) How many bits are required for each entry? 10
c) How many bytes $(1,2,4$ or 8$)$ would you use to store the each entry, so that the size of the FAT is minimal? 2
d) What would be the total size of the FAT in bytes?

$$
2 * 2^{10}
$$

e) How many sectors would the FAT use on the disk?

$$
2^{11} / 2^{10}=2
$$

f) Based on your answers fill in the following table

| Sector number | 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |  |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Content | Boot | FAT | FAT | ROOT | DATA | DATA | DATA | DATA | DATA | D |

g) Assume that FAT [0] contains the start of freelist, and that FileA is being stored on sectors 5,7,4,12 and FileB is being stored on sectors 11,10 . Noting that the last cluster in the file is marked with $\mathbf{0 x F F}$, what would be the content of the FAT table?:

| FAT index 0 | 1 |  | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | ... |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Content | 2 |  | 3 | 6 | 12 | 7 | 8 | 4 | 9 | 13 | FF | 10 | FF | 14 | 15 | 16 | 17 | ... |  |
|  |  |  |  |  |  |  | aKe |  |  |  |  | rec |  |  |  |  |  |  |  |

$\qquad$

