File Abstraction

Most of the following slides are adapted from slides of Randy Bryant of Carnegie Mellon Univ.

UNIX File Abstraction

In UNIX, the file is the basic abstraction used for I/O

Used to access disks, CDs, DVDs, USB and serial devices, network



Unix I/O and C Standard I/O

C Standard

- Most useful for reading/writing files in applications
- Provides buffering between program and actual files

Unix I/O

- Lower level
- Required for system and network programming



Unix I/O Overview

- A Linux *file* is a sequence of *m* bytes:
 - $B_0, B_1, \ldots, B_k, \ldots, B_{m-1}$
- Cool fact: All I/O devices are represented as files:
 - /dev/sda2 (/usr disk partition)
 - /dev/tty2 (terminal)
- Even the kernel is represented as a file:
 - /boot/vmlinuz-3.13.0-55-generic (kernelimage)
 - /proc (kernel data structures)

Unix I/O Overview

- Elegant mapping of files to devices allows kernel to export simple interface called Unix I/O:
 - Opening and closing files
 - open() and close()
 - Reading and writing a file
 - read() and write()
 - Changing the *current file position* (seek)
 - indicates next offset into file to read or write
 - lseek()



Opening Files

Opening a file informs the kernel that you are getting ready to access that file

```
int fd; /* file descriptor */
if ((fd = open("/etc/hosts", O_RDONLY)) < 0) {
    perror("open");
    exit(1);
}</pre>
```

Returns a small identifying integer *file descriptor*

fd == -1 indicates that an error occurred

stdin, stdout, stderr

- In UNIX, every process has three "special" files already open:
 - standard input (stdin) filehandle 0
 - standard output (stdout) filehandle 1
 - standard error (stderr) filehandle 2
- By default, stdin and stdout are connected to the terminal device of the process.
 - Originally, terminals were physically connected to the computer by a serial line
 - These days, we use "virtual terminals" using ssh



VT100 terminal

How the Unix Kernel Represents Open Files

Two descriptors referencing two distinct open disk files. Descriptor 1 (stdout) and 2 (stderr) points to terminal, and descriptor 4 points to file opened on the disk.



File Sharing

- Two distinct descriptors sharing the same disk file through two distinct open file table entries
 - E.g., Calling open twice with the same filename argument



How Processes Share Files: fork()

- A child process inherits its parent's open files
 - Note: situation unchanged by exec() functions
- Before fork() call:



How Processes Share Files: fork()

- A child process inherits its parent's open files
- After fork():
 - Child's table same as parents, and +1 to each refent



Shell redirection

- The shell allows stdin, stdout, and stderr to be redirected (say, to or from a file).
 - > ./myprogram > somefile.txt
 - Connects stdout of "myprogram" to somefile.txt
 - > ./myprogram < input.txt > somefile.txt
 - Connects stdin to input.txt and stdout to somefile.txt
 - > ./myprogram 2> errors.txt
 - Connects stderr to errors.txt
- In this case, the shell simply opens the file, making sure the file handle is 0, 1, or 2, as appropriate.
 - Problem: open() decides what the file handle number is.
 - How do we coerce the filehandle to be 0, 1, or 2?

Initially

stdout prints to the display of the terminal as default.



All we need to do is to point stdout to a file

- Question: But the Descriptor table is kernel space, and we cannot modify it directly.
- Need to use system calls!



dup() : before

#include <unistd.h>
int dup(int filedes);
//dup() returns lowest available file descriptor, now
referring to whatever filedes refers to
newfd = dup(1); // newfd will be 3.



dup() : after

#include <unistd.h>
int dup(int filedes);
//dup() returns lowest available file descriptor, now
referring to whatever filedes refers to
newfd = dup(1); // newfd will be 3.



dup2():before

#include <unistd.h>
int dup2(int oldfd, int newfd);
//Copies descriptor table entry oldfd to entry newfd
int foofd = open("foo.txt", O_WRONLY); //foofd becomes 3.
if (dup2(foofd, stdout)>0) printf("printing to foo.txt\n");



dup2() : after

#include <unistd.h>
int dup2(int oldfd, int newfd);
//Copies descriptor table entry oldfd to entry newfd
int foofd = open("foo.txt", O_WRONLY); //foofd becomes 3.
if (dup2(foofd, stdout)>0) printf("printing to foo.txt\n");



dup() and dup2() pseudocode

dup (fd) returns lowest available file descriptor, now referring to whatever oldfd refers to refers to.

```
//Descriptor table
void *DT[maxFd];
```

```
int dup(int oldfd){
    //get the lowest available
    //file descriptor
    newfd = lowestFd(DT);
    DT(newfd)=DT(oldfd);
    return(newfd);
}
```

dup2(oldfd,newfd) copies descriptor table entry oldfd to entry newfd.

//Descriptor table
void *DT[maxFd];

```
int dup2(int oldfd, int newfd){
    DP[newfd]=DP[oldfd];
    return(newfd);
```

}

- If oldfd is not a valid file descriptor, then the call fails, and newfd is not closed.
- If oldfd is a valid file descriptor, and newfd has the same value as oldfd, then dup2() does nothing, and returns newfd.

I/O and Redirection Example

```
#include "csapp.h"
int main(int argc, char *argv[])
{
    int fd1, fd2, fd3;
   char c1, c2, c3;
   char *fname = argv[1];
   fd1 = open(fname, O RDONLY, 0);
    fd2 = open(fname, O RDONLY, 0);
    fd3 = open(fname, O RDONLY, 0);
   dup2(fd2, fd3);
   read(fd1, &c1, 1);
   read(fd2, &c2, 1);
    read(fd3, &c3, 1);
   printf("c1 = c, c2 = c, c3 = c, c1, c2, c3);
    return 0;
                                             ffiles1.c
```

I/O and Redirection Example



Master Class: Process Control and I/O

```
#include "csapp.h"
int main(int argc, char *argv[])
Ł
    int fd1;
    int s = getpid() \& 0x1;
   char c1, c2;
    char *fname = argv[1];
    fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
    if (fork()) { /* Parent */
        sleep(s);
        Read(fd1, &c2, 1);
       printf("Parent: c1 = %c, c2 = %c n", c1, c2);
    } else { /* Child */
        sleep(1-s);
        Read(fd1, &c2, 1);
        printf("Child: c1 = %c, c2 = %c n'', c1, c2);
    }
    return 0;
                                            ffiles2.c
```

Master Class: Process Control and I/O

```
#include "csapp.h"
                                       Child: c1 = a, c2 = b
int main(int argc, char *argv[])
                                       Parent: c1 = a, c2 = c
Ł
   int fd1;
   int s = getpid() & 0x1;
   char c1, c2;
                                       Parent: c1 = a, c2 = b
   char *fname = argv[1];
                                       Child: c1 = a, c2 = c
   fd1 = Open(fname, O RDONLY, 0);
   Read(fd1, &c1, 1);
   if (fork()) { /* Parent */
                                       Bonus: Which way does it go?
       sleep(s);
       Read(fd1, &c2, 1);
       printf("Parent: c1 = %c, c2 = %c n", c1, c2);
    } else { /* Child */
       sleep(1-s);
       Read(fd1, &c2, 1);
       printf("Child: c1 = %c, c2 = %c n", c1, c2);
    }
   return 0;
                                          ffiles2.c
```

For Further Information

The Unix bible:

- W. Richard Stevens & Stephen A. Rago, Advanced Programming in the Unix Environment, 2nd Edition, Addison Wesley, 2005
 - Updated from Stevens' 1993 book

Stevens is arguably the best technical writer ever.

- Produced authoritative works in:
 - Unix programming
 - TCP/IP (the protocol that makes the Internet work)
 - Unix network programming
 - Unix IPC programming

https://github.com/shihyu/Linux_Programming/tree/master/books

Bonus material

- The following slides are provided as extra and is not part of the course coverage.
- Enjoy!

System Call Error Handling

- On error, Linux system-level functions typically return -1 and set global variable errno to indicate cause.
- Hard and fast rule:
 - You must check the return status of every system-level function
 - Only exception is the handful of functions that return void

Example:

```
if ((pid = fork()) < 0) {
    fprintf(stderr, "fork error: %s\n", strerror(errno));
    exit(-1);
}</pre>
```

Error-reporting functions

Can simplify somewhat using an *error-reporting function*:



Error-handling Wrappers

We simplify the code we present to you even further by using Stevens-style error-handling wrappers:

```
pid_t Fork(void)
{
    pid_t pid;
    if ((pid = fork()) < 0)
        unix_error("Fork error");
        return pid;
}</pre>
```

pid = Fork();

NOT what you generally want to do in a real application

Standard I/O Streams

- Standard I/O models open files as streams
 - Abstraction for a file descriptor and a buffer in memory.
- C programs begin life with three open streams (defined in stdio.h)
 - stdin (standard input)

- stdout (standard output)
- stderr (standard error)

```
#include <stdio.h>
extern FILE *stdin; /* standard input (descriptor 0) */
extern FILE *stdout; /* standard output (descriptor 1) */
extern FILE *stderr; /* standard error (descriptor 2) */
int main() {
   fprintf(stdout, "Hello, world\n");
}
```

Buffering in Standard I/O

Standard I/O functions use buffered I/O



write(1, buf, 6);

Buffer flushed to output fd on "\n" or fflush() call

Standard I/O Buffering in Action

You can see this buffering in action for yourself, using the always fascinating Unix strace program:

```
#include <stdio.h>
int main()
{
    printf("h");
    printf("e");
    printf("l");
    printf("l");
    printf("l");
    printf("o");
    printf("\n");
    fflush(stdout);
    exit(0);
}
```

```
linux> strace ./hello
execve("./hello", ["hello"], [/* ... */]).
...
write(1, "hello\n", 6...) = 6
...
_exit(0) = ?
```

strace: a debugging tool in Linux. When you start a program using **strace**, it prints a list of system calls made by the program.

Fork Example #2 (Earlier Lecture)

- Key Points
 - Both parent and child can continue forking



		Bye
	L1	Bye
		Вуе
L0	L1	Bye

Fork Example #2 (modified)

Removed the "\n" from the first printf

As a result, "L0" gets printed twice



Bye

Bye

Bye

Bye

Repeated Slide: Reading Files

Reading a file copies bytes from the current file position to memory, and then updates file position

Returns number of bytes read from file fd into buf

- Return type ssize_t is signed integer
- **nbytes** < 0 indicates that an error occurred</p>
- short counts (nbytes < sizeof(buf)) are possible and are not errors!

Dealing with Short Counts

Short counts can occur in these situations:

- Encountering (end-of-file) EOF on reads
- Reading text lines from a terminal
- Reading and writing network sockets or Unix pipes

Short counts never occur in these situations:

- Reading from disk files (except for EOF)
- Writing to disk files